



The O-GEHL branch predictor

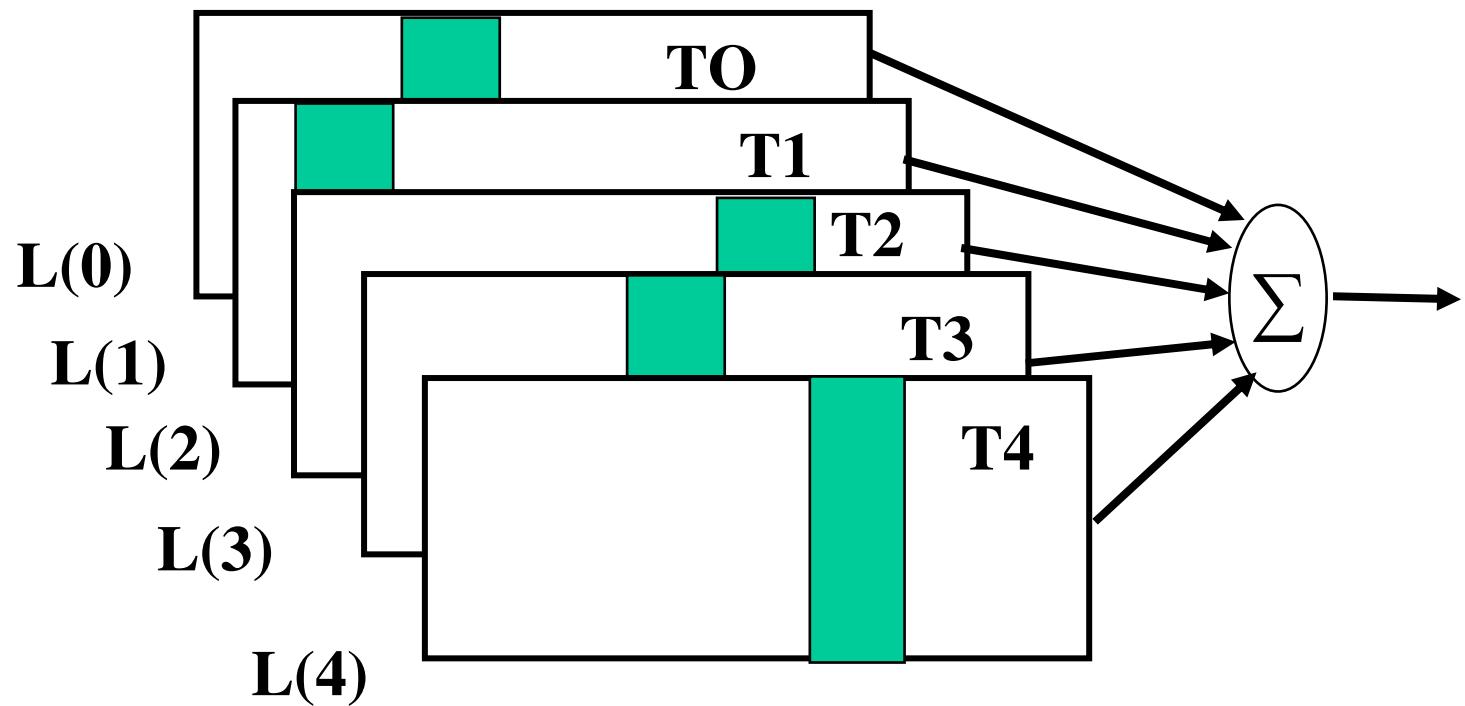
Optimized GEometric History Length

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What is classic 😊

- Global history based:
 - ➔ Yeh and Patt 91, Pan and So 91
- Use of multiple history lengths:
 - ➔ McFarling 93, Evers et al. 96
- Use an adder tree instead of a meta-predictor
 - ➔ Vintan and Iridon 99, Jiménez and Lin 01

Multiple history length neural predictor



GEometric History Length predictor

The set of history lengths forms a **geometric** series

$$L(0) = 0$$

$$L(i) = \alpha^{i-1} L(1)$$

$$\{0, 2, 4, 8, 16, 32, 64, 128\}$$

What is important: $L(i)-L(i-1)$ is drastically increasing

Updating the predictor

- Update on misprediction and under a threshold

8-bit counters and perceptron update threshold (29)

→ Would not have qualified for CBP-1 ☹

Dynamic update threshold fitting

On an O-GEHL predictor, best threshold depends on

- the application ☹
- the predictor size ☹
- the counter width ☹

By chance for the best fixed threshold,

updates on mispredictions ≈ updates on correct predictions

Monitor both numbers
and adapt the update threshold

8 components 8 bits counter would qualify for CBP-1 ☺

Counter width on O-GEHL predictors

- 8 bits are just overkilling 😞
- 4 bits are sufficient 😊
- Mixing 5 bits for short histories and 4 bits for long histories is slightly better 😊
- 3 bits are not so bad !!

Adaptative history length fitting

(inspired by Juan et al 98)

($\frac{1}{2}$ applications: $L(7) < 50$)

\neq

($\frac{1}{2}$ applications: $L(7) > 150$)

Let us adapt some history lengths to the behavior of each application

- **8 tables:**
 - T2: $L(2)$ and $L(8)$
 - T4: $L(4)$ and $L(9)$
 - T6: $L(6)$ and $L(10)$

Adaptative history length fitting (2)

Intuition:

- if high degree of aliasing on T7, stick with short history

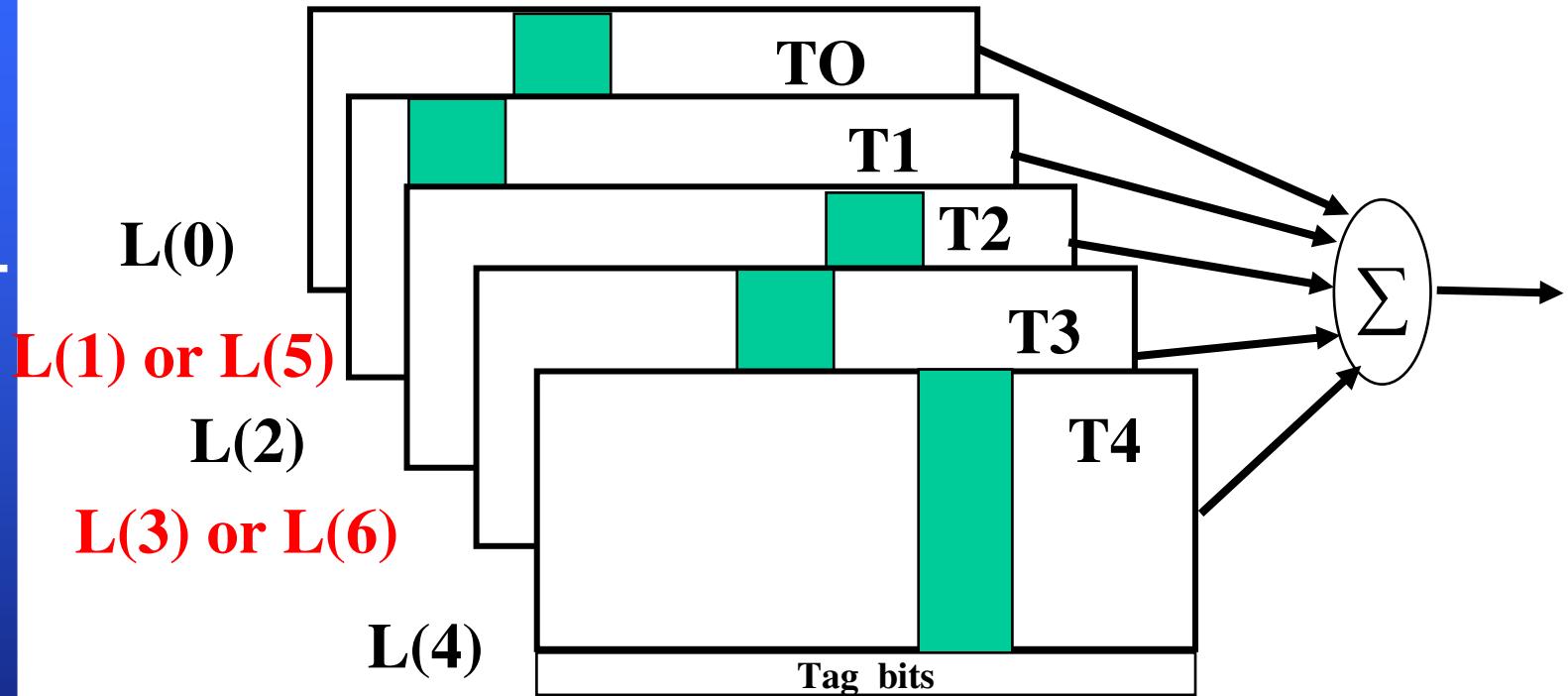
Implementation:

- monitoring of aliasing on updates on T7 through a tag bit and a counter

Simple is sufficient:

- Flipping from short to long histories and vice-versa

The O-GEHL branch predictor



Information to be hashed

- Address + conditional branch history:
 - path confusion on short histories ☹
 - Address + path:
 - Direct hashing leads to path confusion ☹
-
1. Represent all branches in branch history
 2. Use also path history (1 bit per branch, limited to 16 bits)

Configuration for CBP

- 8 tables:
 - ➔ 2 Kentries except T1, 1Kentries
 - ➔ 5 bit counters for T0 and T1, 4 bit counters otherwise
 - ➔ 1 Kbits of one bit tags associated with T7
$$10K + 5K + 6 \times 8K + 1K = 64K$$
- $L(1) = 3$ and $L(10) = 200$
 - ➔ $\{0, 3, 5, 8, 12, 19, 31, 49, 75, 125, 200\}$

Hashing 200+ bits for indexing !!

- Need to compute 11 bits indexes :
 - ➔ Full hashing is unrealistic
- 1. Just regularly pick at most 33 bits in:
address+branch history +path history
- 2. A single 3-entry exclusive-OR stage

A case for the OGEHL predictor (1)

- High accuracy
- Robustness to variations of history lengths choices:
 - ➔ L(1) in [2,6], L(10) in [125,300]
 - ➔ misp. rate < 1.04 x reference misp. rate
- Geometric series: not a bad formula !!
 - ➔ best geometric L(1)=3, L(10)=223, REF-0.02 misp/KI
 - ➔ best overall {0, 2, 4, 9, 12, 18, 31, 54, 114, 145, 266} REF-0.04 misp/KI

A case for the OGEHL predictor (2)

- Reduce counter width by 1 bit: 49 Kbits
 - would have been a finalist ☺
- 64 Kbits 4 components OGEHL predictor
 - would have been a finalist ☺
- 50 Kbits 4 components OGEHL predictor (3-bit)
 - would have been a finalist ☺
- 768 Kbits 12 components OGEHL predictor
 - 2.25 misp/KI

A case for the O-GEHL predictor (3)

- O-GEHL predictor uses only global information
- Can be ahead pipelined
- Prediction computation logic complexity is low

(The End)